

***Computer Program of Line Balancing, Regarding  
Efficiency and Number of Stations as Variables***

Fumio AKAGI\*, Hirokazu OSAKI\*,  
and Susumu KIKUCHI\*

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Synopsis

To assign work elements to the work stations in an assembly or manufacturing line, various computer programs have been developed and used. And it does that the number of stations or the cycle time is even given. But in practice it is desirable to obtain the assignment which shows the highest efficiency of line balancing under all possible combinations of the number of stations and the cycle time.

Therefore we propose a computer program of the assignment method in which the efficiency of line balancing,  $E_{bb}$  and the number of stations,  $NN$  are regarded as variables. In this method the minimum value ( $EE_{bb}$ ) of efficiency and the constant term ( $d$ ) by which  $E_{bb}$  is reduced are given previously. And for any combination of  $E_{bb}$  ( $EE_{bb} \leq E_{bb} \leq 100$ ) and  $NN$  ( $1 \leq NN \leq N_m : N_m$  calculated from  $EE_{bb}$ ), the work elements are assigned to work stations, the precedence restrictions being used.  $E_{bb}$  is reduced by  $d$  from the ideal value (100) until the assignment to  $NN$  is obtained. The efficiency of the obtained assignment,  $E_{bo}$  is calculated. As  $E_{bb} < E_{bo} \leq E_{bb} + d$ , the calculation is continued until the assignment to  $NN$ , which shows the maximum efficiency, is obtained. In this process  $NN$  varies from  $N_m$  to 1 by 1.

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\* Department of Industrial Science, School of Engineering  
Okayama University.

## 1. Introduction

The basic problem in line balancing is to assign the work elements to the work stations so as to balance the workload (the station time) at each station and to make this work content as close to the cycle time of the line as possible [1]. Various computer programs have been developed and used to solve this problem [2,3]. Generally to evaluate the assignment, the efficiency of line balancing, which is indicated by the following  $E_b$ , has been used [4].

$$E_b = \frac{T}{N \times P} \times 100 \quad (1)$$

where  $E_b$  = the efficiency of line balancing

$T$  = the total work time

$N$  = the number of stations

$P$  = the cycle time to  $N$  stations.

As there is no slack time in the line at  $E_b=100\%$ , the assignment of the work elements is ideal.

In the programs developed until now,  $N$  or  $P$  must be given. As  $N$  is given, the assignment is determined so as to minimize  $P$ . As  $P$  is given, it is done so as to minimize  $N$ . The number of stations  $N$  is determined by personel, machines, tools, work areas, and so on. The cycle time  $P$  is restricted by the production time and the units to be produced per day. From these view points the number of stations ( $NN$ ), the efficiency of line balancing ( $E_{bb}$ ), and the cycle time ( $PP$ ) are regarded as variables in some restrictions. Further there is a functional relation among  $NN$ ,  $PP$  and  $E_{bb}$ .

Therefore in this paper we propose an assignment method in which the work elements are assigned to the work stations,  $E_{bb}$  and  $NN$  regarded as variables. Firstly, in this method the minimum efficiency of line balancing ( $EE_b$ ) and the constant term ( $d$ ) by which  $E_{bb}$  is reduced are given. Then the maximum number of stations ( $N_m$ ) is determined by  $EE_b$  and the maximum work element time ( $t_m$ ) from Eq.(1), and the value of  $NN$  is given as  $N_m$ . Next, the value of  $E_{bb}$  is given as 100 (the ideal value). Then  $PP$  is determined from  $E_{bb}$  and  $NN$  from Eq. (1).  $NN$  and  $PP$  being used, the work elements are assigned to the work stations according to the precedence restrictions so that the work time of each station (the station time) may be less than  $PP$ .  $E_{bb}$  is reduced at the constant rate ( $d$ ) from 100 until the assignment to  $NN$  is obtained. And the efficiency of the obtained assignment ( $E_{bo}$ ) is calculated.  $E_{bo}$  is greater than  $E_{bb}$  but less than  $E_{bb}+d$  because the maximum station time ( $C$ ) is less than  $PP$ . Since then  $E_b$  is made

exchange Ebb to Ebo.

And the calculation is continued until the assignment to NN which shows the maximum of Eb between Ebb and Ebb+d is obtained. The same procedures are repeated from Nm to 1 by 1 by 1.

As this method needs much repetition of the routine calculation, we develop the computer program.

## 2. Procedures for Assigning Work Elements to Work Stations

In order to obtain a favorable assignment, the sequence of the operations must be analyzed first. The precedence restrictions are used to indicate which work element must be done before others. The results are summarized by a precedence diagram[5], an arrow diagram[6], a table of functional units[7], a precedence matrix[1], and so on.

The proposed procedure for assigning the work elements to the work stations using the precedence restrictions is as follows.

Step 1. EEb and d are given.

Step 2. The theoretical maximum number of stations, Nm is calculated from

$$N_m = \left[ \frac{T}{t_m} \times \frac{100}{EEb} \right] \quad (2)$$

where  $t_i$  = the operation time of work element  $i$

$t_m$  = max  $t_i$ , the maximum work element time

$T$  =  $\sum t_i$ , the total work time

[ ] : Gaussian symbol.

Therefore the possible range of NN is from Nm to 1, that is,  $1 \leq NN \leq N_m$ . For any NN, the following procedures are applied.

Step 3. At the first time NN is given as Nm. After the second time NN is down from Nm by 1, that is,  $NN = NN - 1$ . Go to next step.

When NN equals 1, stop the procedure.

Step 4. PP to NN and Ebb is calculated by

$$PP = \frac{T}{NN} \times \frac{100}{Ebb} \quad (3)$$

Ebb is reduced by d from 100 until an assignment to NN is obtained. To any combination (NN, PP), go to next step.

Step 5. At the first time, select the work elements which don't have the precedence elements.

After the second time, add to them the work elements which have only the assigned ones as the precedence ones. Go to next step.

Step 6. Select the work elements, which is less than the remaining cycle time (the slack time) of the station, from among the work elements

selected at step 5. Then go to step 7.

If no work element can be selected here, proceed to the next station and repeat the same step.

If the last station has been examined, go to step 8.

Step 7. Assign the work element having the maximum work time in those elements selected at step 6 to the station. Subtract the work time from the remaining cycle time and eliminate the assigned work element. Then return step 5.

Step 8. When all the elements has been assigned to all the stations, an assignment has been obtained.

The efficiency of the obtained assignment ( $E_{bo}$ ) is calculated from

$$E_{bo} = \frac{T}{NN \times C} \times 100 \quad (4)$$

where  $C$  = the maximum station time of the obtained assignment for  $NN$  and  $PP$

$$C < PP \quad .$$

Because of  $E_{bb} < E_{bo} \leq E_{bb} + d$ , since then we search the assignment which shows the maximum efficiency. That is, let  $E_{bb} = E_{bo}$  and  $PP = C$ . To the new combination of  $(NN, PP)$ , return step 5.

If one or more work elements remain and an assignment to  $NN$  hadn't been obtained, return step 4 and reduce  $E_{bb}$  by  $d$ , that is,  $E_{bb} = E_{bb} - d$ .

If one or more work elements remain and an assignment to  $NN$  had been ever obtained, return step 3 and reduce  $NN$  by 1.

### 3. Program

This program is written in Fortran IV and is the form of the subroutine.

The subroutine name is CPLB.

```
SUBROUTINE CPLB(EEB,D,NWORK,NAME,TIME,NAO,KINDP,NAMEP,NSTART)
```

#### 3.1. Usage

The work time and the precedence restrictions among work elements are provided from the table of functional units or the arrow diagram.

##### 3.1.1 Argument List in the case of Table of Functional Units

<u>ARGUMENT</u>	<u>I/O</u>	<u>TYPE</u>	<u>SIZE</u>	<u>DEFINITION</u>
EEB	I	REAL	1	$E_{bb}$ , the minimum efficiency of line balancing
D	I	REAL	1	$d$ , the constant term by which $E_{bb}$ is reduced
NWORK	I	INTEGER	1	number of work elements
NAME	I	CHARACTER	50X5	name of work element (5A4)

<u>ARGUMENT</u>	<u>I/O</u>	<u>TYPE</u>	<u>SIZE</u>	<u>DEFINITION</u>
TIME	I	REAL	50	$t_i$ , time of work element
NAO	I	INTEGER	50	work element's number
KINDP	I	INTEGER	50	number of parts in work element
NAMEP	I	INTEGER	50	part's number in work element
NSTART	I	INTEGER	1	0

where  $C < D < EEB \leq 100$ ,  $NWORK \leq 50$ ,  $NAME \leq 20$  characters,  $KINDP \leq 10$

### 3.1.2 Argument List in the case of Arrow Diagram

<u>ARGUMENT</u>	<u>I/O</u>	<u>TYPE</u>	<u>SIZE</u>	<u>DEFINITION</u>
EEB	I	REAL	1	EEb, the minimum efficiency of line balancing
D	I	REAL	1	d, the constant term by which Ebb is reduced
NWORK	I	INTEGER	1	number of work elements
NAME	I	CHARACTER	50X5	name of work element (5A4)
TIME	I	REAL	50	$t_i$ , time of work element
NAO	I	INTEGER	50	preceding event number
KINDP	I	INTEGER	50	succeeding event number
NAMEP	I	INTEGER	50	not use in this case
NSTART	I	INTEGER	1	source node number

where  $0 < D < EEB \leq 100$ ,  $NWORK \leq 50$ ,  $NAME \leq 20$  characters.

### 3.2. Suggestion on Using

Subroutine GRPUNT or GRPARR, PRINT, and MAXGRP are used in CPLB. GRPUNT is used to select work elements which don't have the precedence works using the table of functional units. GRPARR to do so using the arrow diagram. PRINT to print out the results of the assignment. And MAXGRP to select the work element having the maximum work time from the selected work elements.

Program list is shown in Table 1.

## 4. Example

The assembly work of a small electric switch is used as an example to excute the program. The work has been analyzed and divided into work elements. The table of the functional units and the list of the arrow diagram have been developed.

Either the table or the list is used as the input data of the precedence restrictions among the work elements. The computer outputs are the same.

The data given from the table of the functional units are shown in Table 2. And the data given from the list of the arrow diagram are shown in Table 3. The computer outputs of this example are shown in Table 4.

#### References

- [1] ASTM : "Manufacturing Planning and Estimating Handbook", McGraw-Hill(1963), 9-41.
- [2] K. Takeda : IE, 5(1975), 39.
- [3] E. M. Darel : AIIE Trans., 7-3(1975), 302.
- [4] M. Kawashima : "Shin Sagyo to Layout", Nihon Noritsu Kyokai(1964) 199.
- [5] T. O. Prenting and R. M. Battglin : J. Indust. Engng., 15-4(1964-7), 208.
- [6] K. Tone : "PERT Nyumon", Toyo keizai(1973), 18.
- [7] F. Akagi, H. Osaki and S. Kikuchi : Conference Report of J. S. M. E., 780-8(1978-4), 24.

Table 1. Program Listing

```

SUBROUTINE CPLB(EEB,D,NWORK,NAME,TIME,NAO,KINDP,NAMEP,NSTART)
DIMENSION NAME(50,5),TIME(50),NAO(50),KINDP(51),NAMEP(50,10)
DIMENSION NGROUP(50),NSGRP(50),NSGRP1(50),NAOS(50),NAOSS(50)
DIMENSION CYCLET(50),NKELMT(50),NSTELM(50,50)
BIG=0.0
SUM=0.0
DO 100 I=1,NWORK
SUM=SUM+TIME(I)
IF(TIME(I) .LE. BIG) GO TO 100
BIG=TIME(I)
100 CONTINUE
STN=SUM/EEB/BIG*100.0
NSTN=STN
90 CONTINUE
IF(NSTN .LE. 0) GO TO 91
EB=1.0
KDOWN=0
TMEAN=SUM/FLGAT(NSTN)
55 CONTINUE
TMAX=TMEAN/EB
50 CONTINUE
TMIN=TMEAN*(2.0-1.0/EB)
DO 101 I=1,NWORK
NAOS(I)=NAO(I)
NAOSS(I)=NAO(I)
101 CONTINUE
DO 110 I=1,NSTN
CYCLET(I)=0.0
NKELMT(I)=0
110 CONTINUE
K=0
NUMGRP=0
NPLACE=0
MAXET=51
KINDP(51)=NSTART
10 CONTINUE
K=K+1
IF(K .GT. NSTN) GO TO 30
L=0
20 CONTINUE
IF(NSTART .LE. 0) GO TO 1
CALL GRPARR(NWORK,NUMGRP,NGROUP,NAOS,NAOSS,KINDP,NAO,MAXET)
GO TO 2
1 CONTINUE
CALL GRPUNT(NWORK,NUMGRP,NGROUP,NAOS,NAOSS,KINDP,NAMEP)
2 CONTINUE
DO 250 I=1,NUMGRP
NUM=NGROUP(I)
IF(NAOS(NUM) .LE. 0) GO TO 250
IF(TIME(NUM) .LE. 0.0) GO TO 61
250 CONTINUE
K1=0

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TREST=TMAX-CYCLET(K)
NTREST=TREST*10000.0
TREST=NTREST/10000
DO 300 I=1,NUMGRP
  K2=NGROUP(I)
  IF(NAOS(K2) .EQ. 0) GO TO 300
  IF(TREST .LE. TIME(K2)) GO TO 300
  K1=K1+1
  NSGRP(K1)=NGROUP(I)
300 CONTINUE
  IF(K1 .EQ. 0) GO TO 10
  MAXET=MAXGRP(K1,NSGRP,TIME)
  CYCLET(K)=CYCLET(K)+TIME(MAXET)
  NKELMT(K)=NKELMT(K)+1
  L=NKELMT(K)
  NSTELM(K,L)=MAXET
  GO TO 60
61 MAXET=NUM
60 CONTINUE
  NAOS(MAXET)=0
  NPLACE=NPLACE+1
  GO TO 20
30 CONTINUE
  IF(NPLACE .GE. NWORK) GO TO 41
  IF(KDOWN .NE. 0) GO TO 40
  EB=EB-(D/100.0)
  GO TO 55
41 CONTINUE
  CALL PRINT(SUM,NSTN,CYCLET,NKELMT,NSTELM,NAO,NAME,TIME,EB,TMAX,
&          TMIN,TMEAN)
  KDOWN=1
  GO TO 50
40 CONTINUE
  WRITE(6,6800)
6800 FORMAT(////,20X,'DIVIDE THE WORK ELEMENT OR REDUCE THE STATION',)
  NSTN=NSTN-1
  GO TO 90
91 CONTINUE
  RETURN
  END

```

```

C ** SELECT ASSIGNABLE WORK ELEMENT FROM TABLE OF FUNCTIONAL UNITS **
SUBROUTINE GRPUNT(NWORK,NUMGRP,NGROUP,NAOS,NAOSS,KINDP,NAMEP)
DIMENSION NGROUP(50),NAOS(50),NAOSS(50),KINDP(51),NAMEP(50,10)
DO 200 I=1,NWORK
  IF(NAOSS(I) .EQ. 0) GO TO 200
  KP=KINDP(I)
  DO 210 J=1,KP
    DO 211 I1=1,NWORK
      IF(NAOS(I1) .EQ. 0) GO TO 211
      IF(NAMEP(I,J) .EQ. NAOS(I1)) GO TO 200
211 CONTINUE
210 CONTINUE
  NUMGRP=NUMGRP+1
  NGROUP(NUMGRP)=I
  NAOSS(I)=0
200 CONTINUE
  RETURN
  END

```



```

C ** SELECT ASSIGNABLE WORK ELEMENT FROM ARROW DIAGRAM **
SUBROUTINE GRPARR(NWORK,NUMGRP,NGROUP,NAOS,NAOSS,NBACK,NFRONT,
&
      MAXET)
  DIMENSION NGROUP(50),NAOS(50),NAOSS(50),NBACK(51),NFRONT(50)
  DO 100 I=1,NWORK
    IF(NAOS(I) .LE. 0) GO TO 100
    IF(NBACK(I) .EQ. NBACK(MAXET)) GO TO 300
100  CONTINUE
    DO 200 I=1,NWORK
      IF(NAOSS(I) .LE. 0) GO TO 200
      IF(NFRONT(I) .NE. NBACK(MAXET)) GO TO 200
      NUMGRP=NUMGRP+1
      NGRUP(NUMGRP)=I
      NAOSS(I)=0
200  CONTINUE
300  CONTINUE
    RETURN
    END

C ** PRINT OUT THE SOLUTION OF ASSIGNMENT **
SUBROUTINE PRINT(SUM,NSTN,CYCLET,NKELMT,NSTELM,
&
      NAO,NAME,TIME,EB,TMAX,TMIN,TMEAN)
  DIMENSION CYCLET(50),NKELMT(50),NSTELM(50,50),NAO(50),
&
      NAME(50,10),TIME(50)
  WRITE(6,6100) NSTN,EB,SUM,TMEAN,TMAX,TMIN
6100 FORMAT(1H ,///,30X,'THE INITIAL DATA',/,10X,'NUM. OF STATION =',
  1      15,5X,'EB=',F5.3,5X,'SUM=',F10.5,2X,'TMEAN=',F10.5,2X,
  2      'TMAX=',F10.5,2X,'TMIN=',F10.5,/,10X,'STATION',3X,
  3      'CYCLE TIME',3X,'WORK ELEMENT',/)
  BIG=0.0
  DO 600 I=1,NSTN
    WRITE(6,6200) I,CYCLET(I),NKELMT(I)
6200  FORMAT(10X,15,3X,F10.5,15)
    L=NKELMT(I)
    DO 610 J=1,L
      J2=NSTELM(I,J)
      WRITE(6,6300) NAO(J2),(NAME(J2,J3),J3=1,4),TIME(J2)
6300  FORMAT(30X,15,5X,4A4,F10.4)
610  CONTINUE
    IF(CYCLET(I) .LE. BIG) GO TO 600
    BIG=CYCLET(I)
600  CONTINUE
    EB=SUM/(NSTN*BIG)
    AG=FLOAT(NSTN)*BIG/SUM-1.0
    WRITE(6,6400) BIG,EB,AG
6400  FORMAT(//,20X,'TMAX=',F10.5,5X,'EB=',F5.3,5X,'AG=',F5.3)
    TMAX=BIG
    RETURN
    END

FUNCTION MAXGRP(NUMBER,NSET,DATA)
  DIMENSION NSET(50),DATA(50)
  BIG=0.0
  DO 100 I=1,NUMBER
    NUM=NSET(I)
    IF(DATA(NUM) .LE. BIG) GO TO 100
    MAXGRP=NUM
    BIG=DATA(NUM)
100  CONTINUE
  RETURN
  END

```

Table 2. Given Data from Table of Functional Units

EEB = 80.0

D = 5.0

NWORK= 6

<u>NAME (i,j)</u>	<u>TIME (i)</u>	<u>NAO (i)</u>	<u>KINDP (i)</u>	<u>NAMEP (i,k)</u>
BASE	32.3	1	7	101 802 803 804 805 306 407
GUIDE	15.3	2	4	802 408 309 810
BUTTON	5.8	3	3	811 2 803
TERMINAL 1	12.6	4	3	112 313 805
TERMINAL 2	12.6	5	3	112 313 805
TERMINAL 3	12.6	6	3	112 313 805

NSTART= 0

Table 3. Given Data from List of Arrow Diagram\*\*

EEB = 80.0

D = 5.0

NWORK= 9

<u>NAME (i,j)</u>	<u>TIME (i)</u>	<u>NAO (i)</u>	<u>KINDP (i)</u>
BASE	32.3	1	6
GUIDE	15.3	1	2
BUTTON	5.8	2	6
TERMINAL 1	12.6	1	3
TERMINAL 2	12.6	1	4
TERMINAL 3	12.6	1	5
DUMMY	0.0	3	6
DUMMY	0.0	4	6
DUMMY	0.0	5	6

NSTART= 1

\*\* Arrow Diagram of the Small Electric Switch

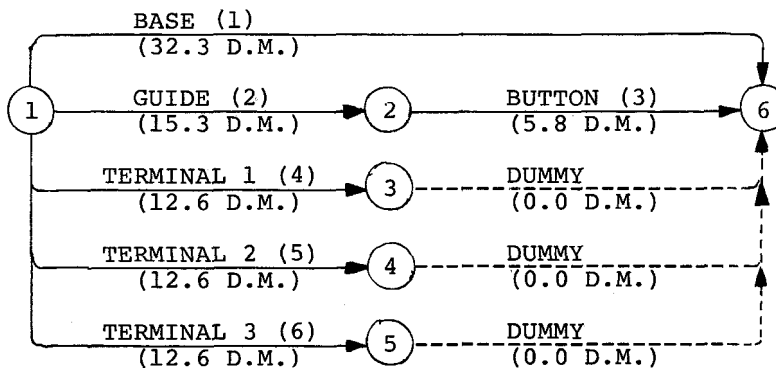


Table 4. Computer Output

NUM. OF STATION =		THE INITIAL DATA		SUM=	TMEAN=
3 * <sup>1</sup>		EB=0.900 * <sup>2</sup>		* <sup>3</sup> 91.20000	30.4000
STATION	CYCLE TIME	WORK ELEMENT		TMAX=	TMIN=
				* <sup>4</sup> 33.77778	27.0222
1	32.30000	1			
		1	BASE	32.3000	* <sup>5</sup>
2	27.90000	2			
		2	GUIDE	15.3000	
		4	TERMINAL 1	12.6000	
3	31.00000	3			
		5	TERMINAL 2	12.6000	
		6	TERMINAL 3	12.6000	
		3	BUTTON	5.8000	

TMAX= 32.30000 \*<sup>6</sup> EB=0.941 \*<sup>7</sup> AG=0.062 \*<sup>8</sup>

NUM. OF STATION =		THE INITIAL DATA		SUM=	TMEAN=
2		EB=0.950		91.20000	45.6000
STATION	CYCLE TIME	WORK ELEMENT		TMAX=	TMIN=
				48.00000	43.2000
1	44.90000	2			
		1	BASE	32.3000	
		4	TERMINAL 1	12.6000	
2	46.30000	4			
		2	GUIDE	15.3000	
		5	TERMINAL 2	12.6000	
		6	TERMINAL 3	12.6000	
		3	BUTTON	5.8000	

TMAX= 46.30000 EB=0.985 AG=0.015

NUM. OF STATION =		THE INITIAL DATA		SUM=	TMEAN=
1		EB=0.950		91.20000	91.2000
STATION	CYCLE TIME	WORK ELEMENT		TMAX=	TMIN=
				96.00000	86.4000
1	91.20000	6			
		1	BASE	32.3000	
		2	GUIDE	15.3000	
		4	TERMINAL 1	12.6000	
		5	TERMINAL 2	12.6000	
		6	TERMINAL 3	12.6000	
		3	BUTTON	5.8000	

TMAX= 91.20000 EB=1.000 AG=\*\*\*\*\*

\*<sup>1</sup> Nm = NN, \*<sup>2</sup> Ebb/100, \*<sup>3</sup> T =  $\sum t_i$ , \*<sup>4</sup> PP, \*<sup>5</sup> tm = max ti,

\*<sup>6</sup> C ,

\*<sup>7</sup> Ebo/100 = T/(NN X C),

\*<sup>8</sup> AG = (NN X C - T) / T .